

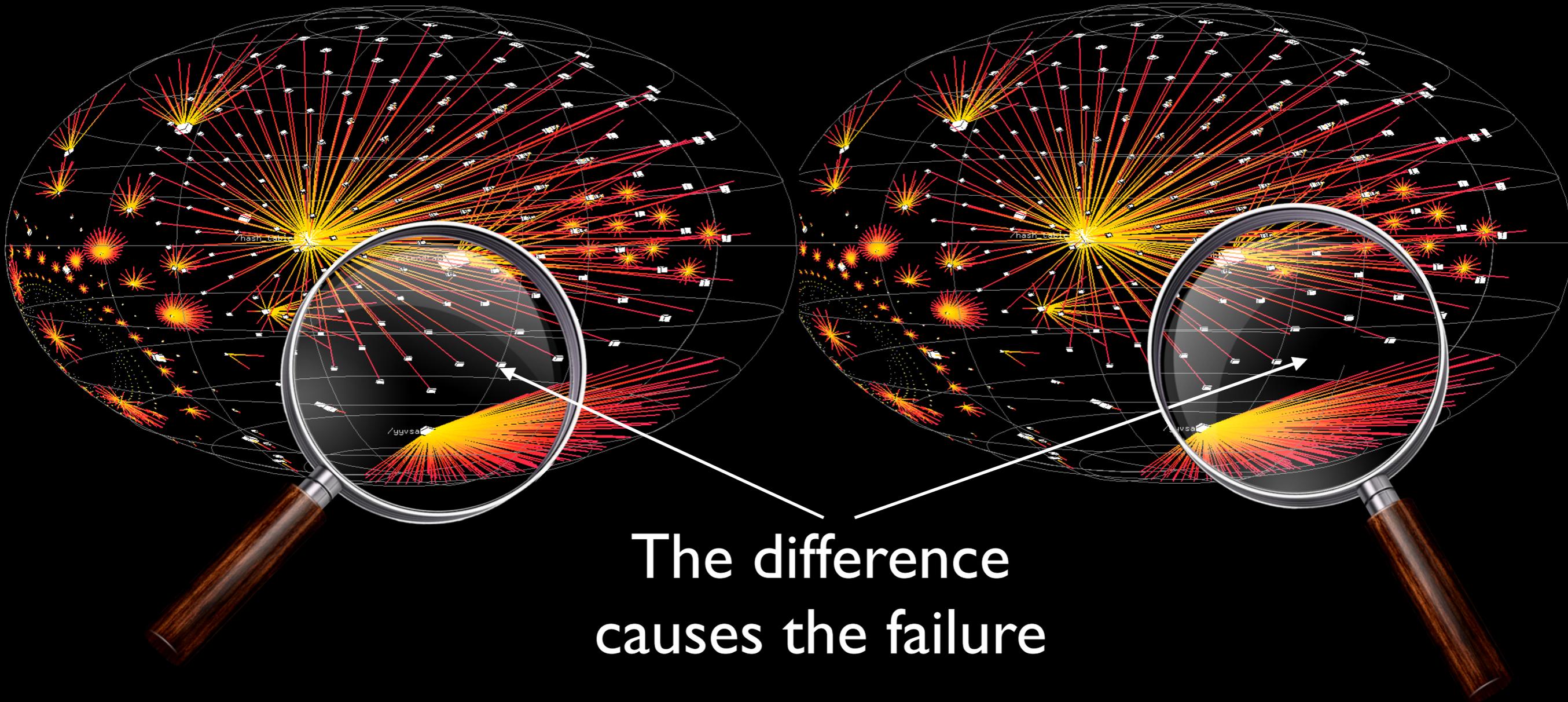
Locating Failure Causes

Andreas Zeller

Finding Causes

Infected state

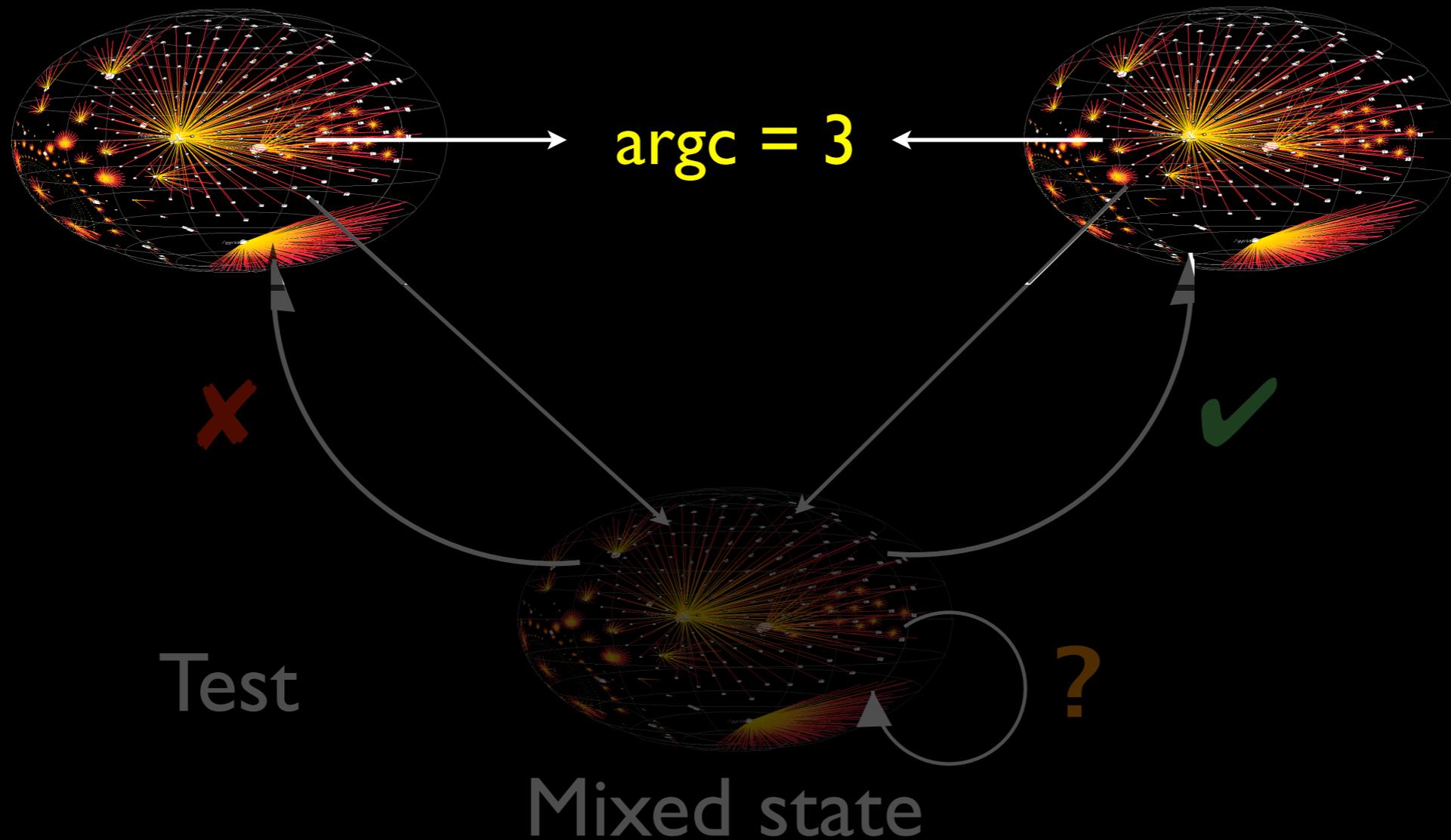
Sane state



Search in Space

Infected state

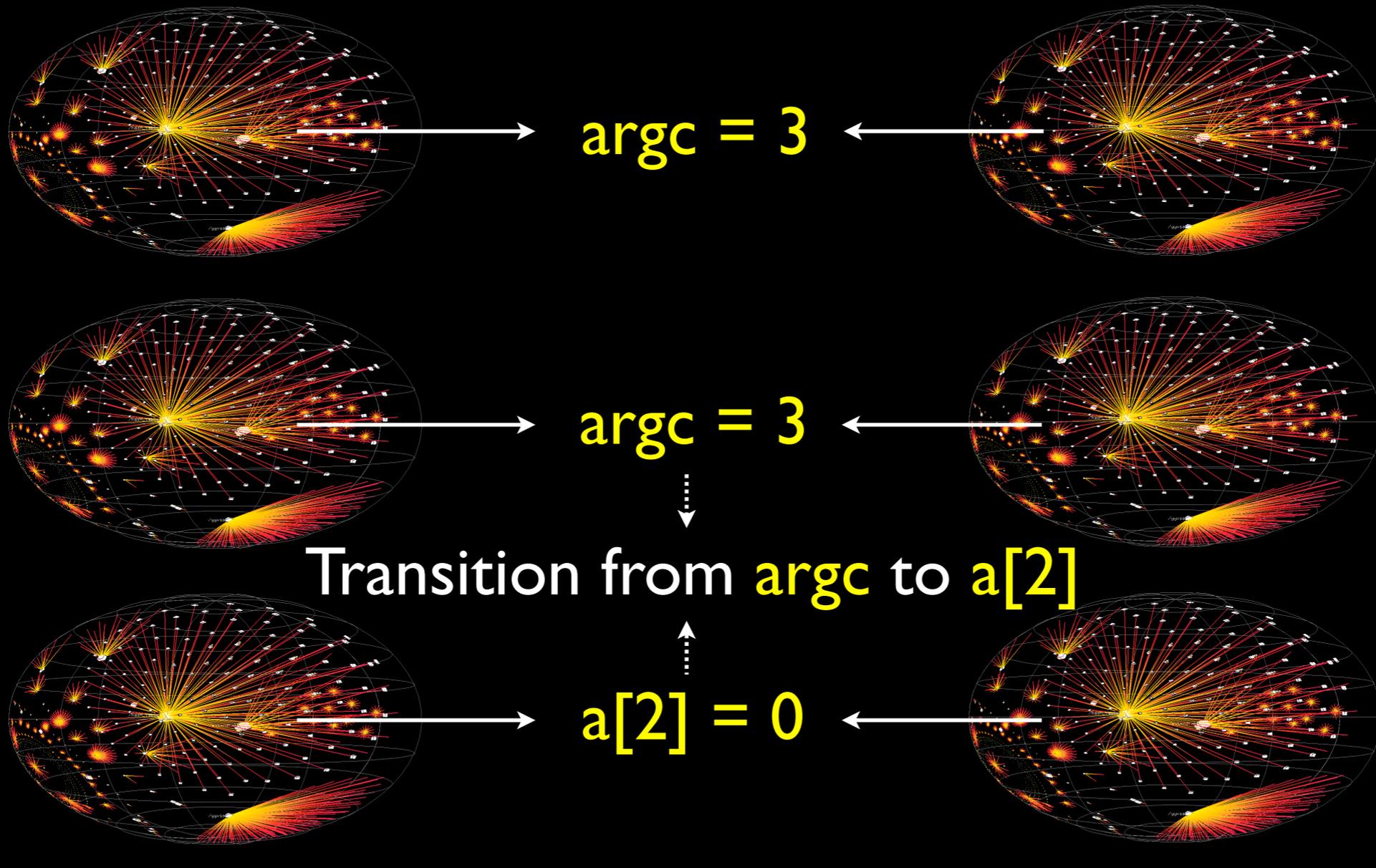
Sane state



Search in Time

Failing run

Passing run



Transitions

A cause transition occurs when a new variable begins to be a failure cause:

- **argc** no longer causes the failure...
- ...but **a[2]** does!

Can be narrowed down by binary search

```
int main(int argc, char *argv[])
{
    int *a;

    // Input array
    a = (int *)malloc((argc - 1) * sizeof(int));
    for (int i = 0; i < argc - 1; i++)
        a[i] = atoi(argv[i + 1]);

    // Sort array
    shell_sort(a, argc);

    // Output array
    printf("Output: ");
    for (int i = 0; i < argc - 1; i++)
        printf("%d ", a[i]);
    printf("\n");

    free(a);
    return 0;
}
```

Should be $argc - 1$

Why Transitions?

- Each failure cause in the program state is caused by some statement
- These statements are executed at **cause transitions**
- Cause transitions thus are **statements that cause the failure!**

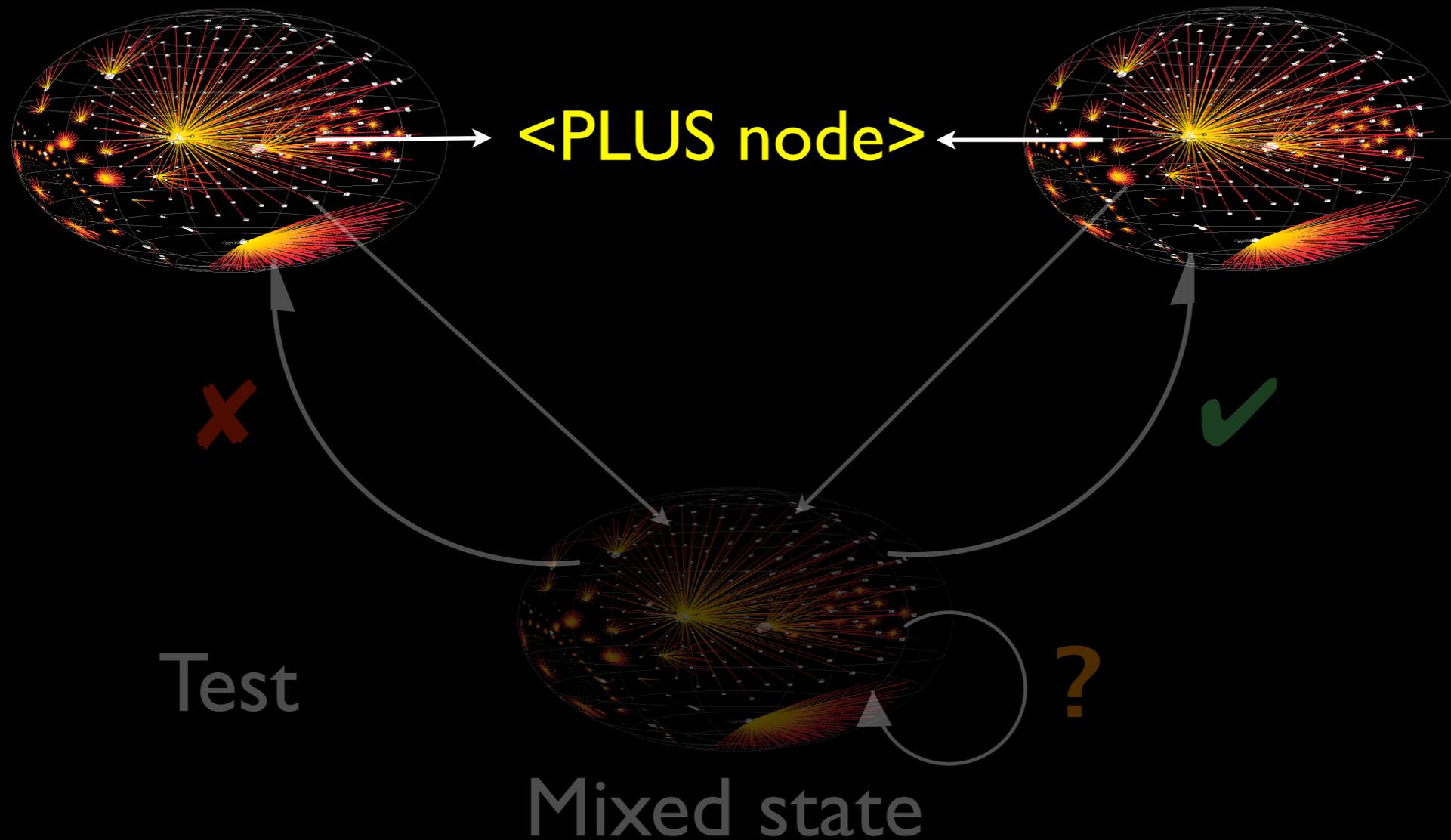
Potential Fixes

- Each cause transition implies a *fix* to make the failure no longer occur – just prohibit the transition
- A cause transition is more than a potential fix – it may be “the” defect itself

Searching GCC State

Infected state

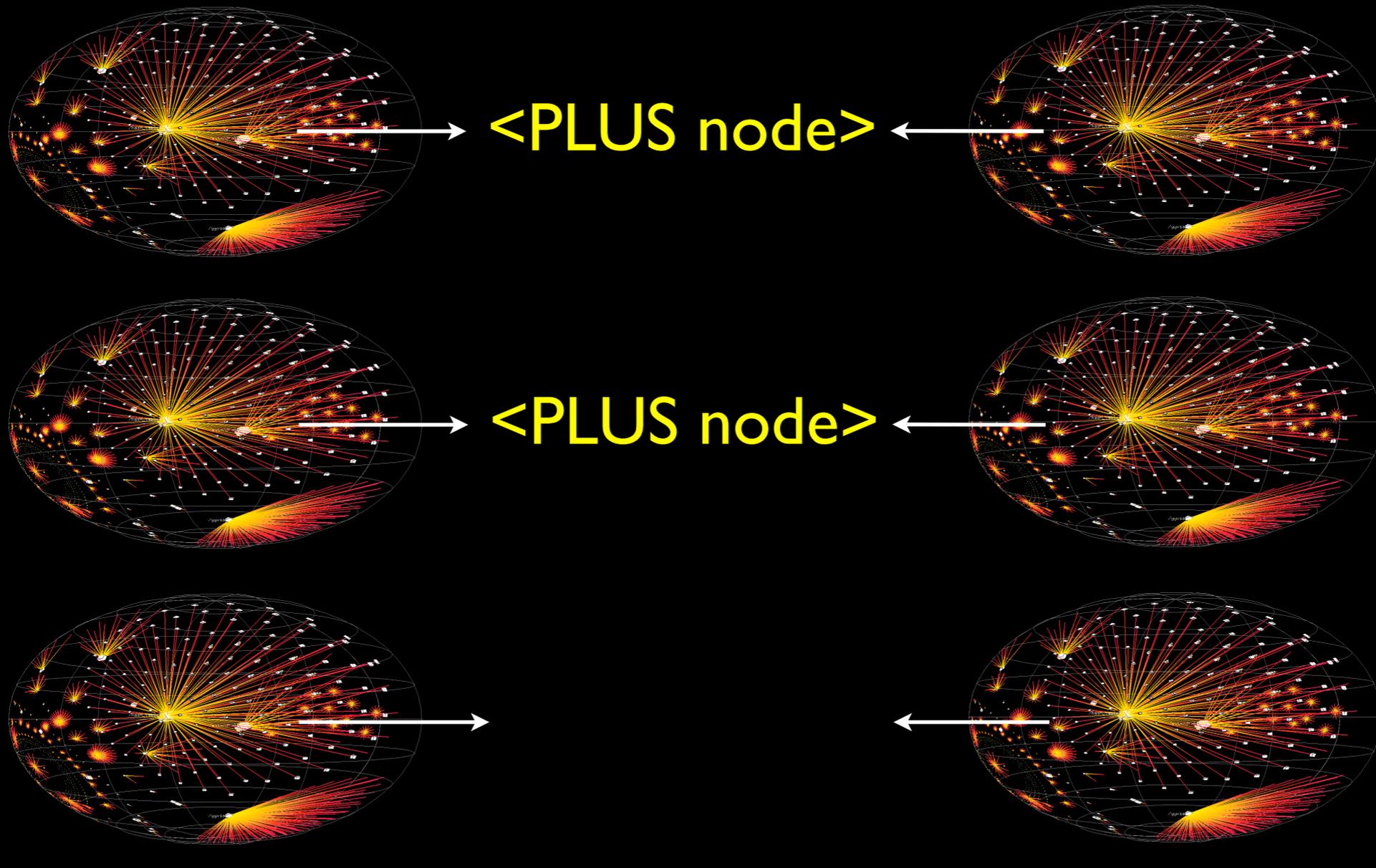
Sane state



Search in Time

Failing run

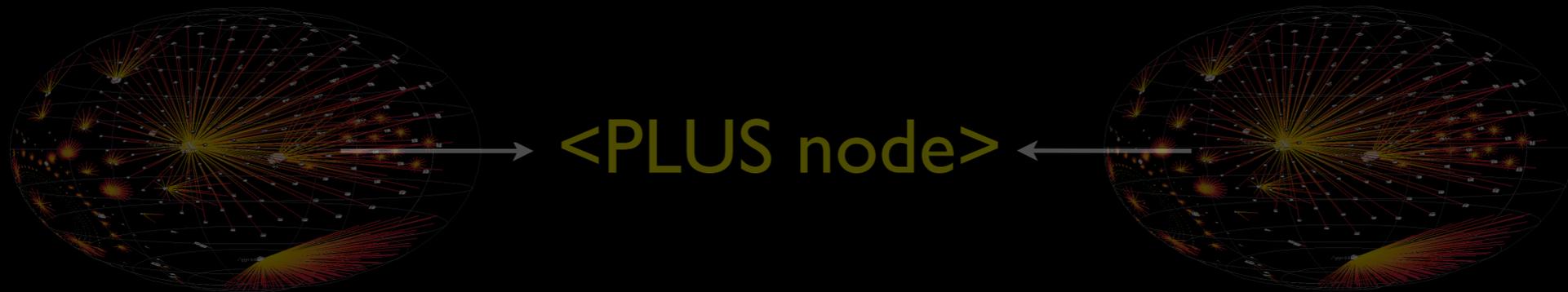
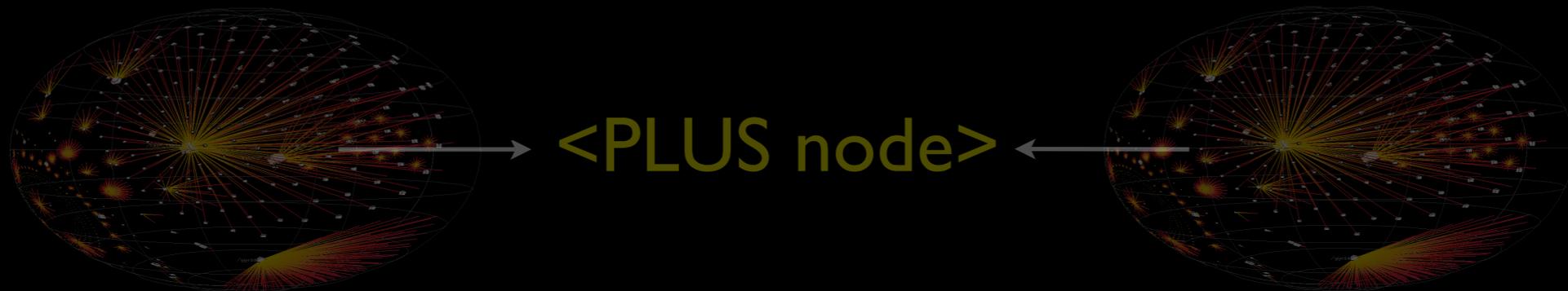
Passing run



Search in Time

Failing run

Passing run



$\text{link} \rightarrow \text{fld}[0].\text{rtx} \rightarrow \text{fld}[0].\text{rtx} == \text{link}$

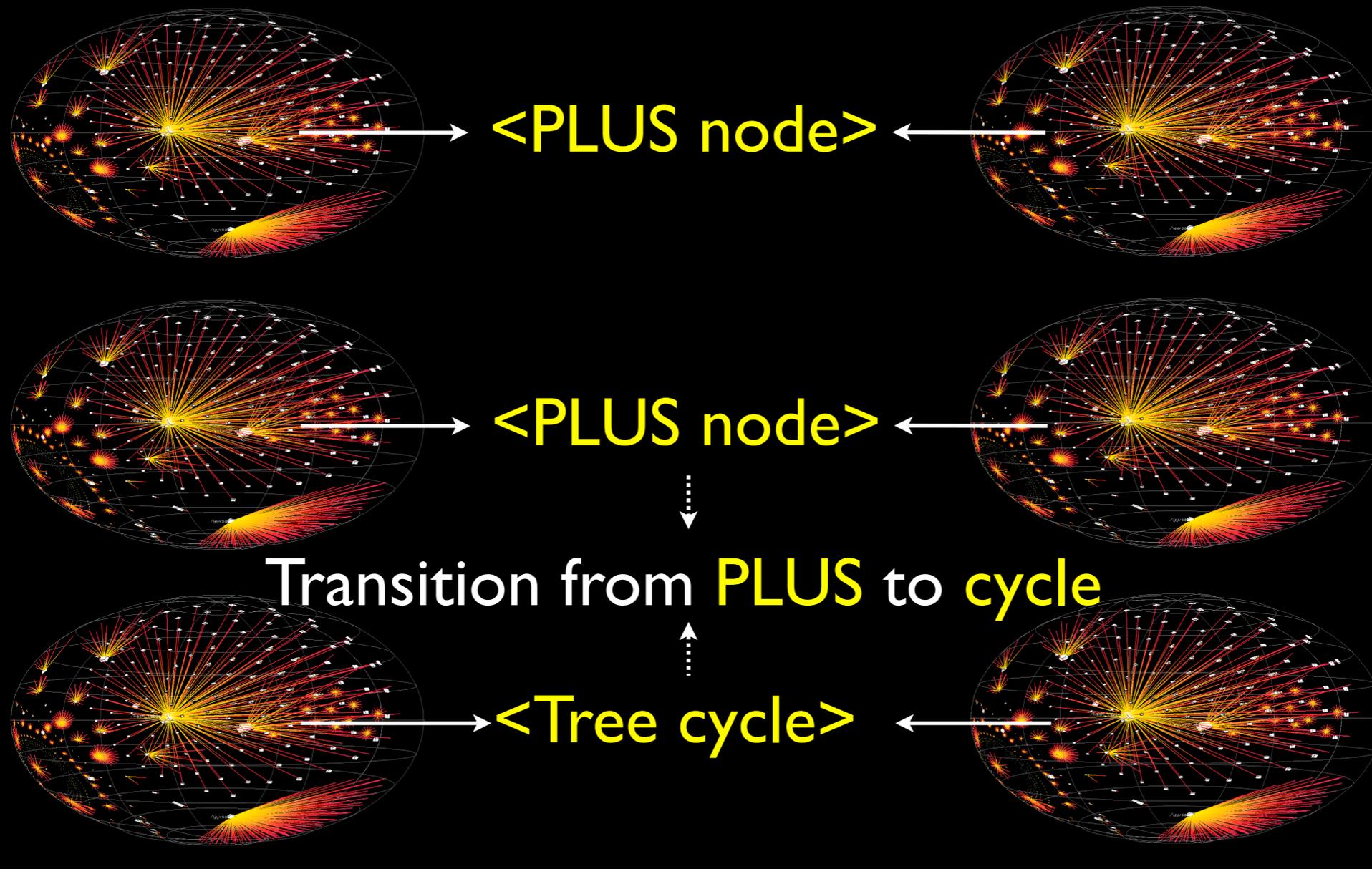


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Search in Time

Failing run

Passing run



All GCC Transitions

#	Location	Cause transition to variable
0	⟨Start⟩	argv[3]
1	tolev.c:4755	name
2	tolev.c:2909	dump_base_name
3	c-lex.c:187	finput→_IO_buf_base
4	c-lex.c:1213	nextchar
5	c-lex.c:1213	yyssa[41]
6	c-typeck.c:3615	yyssa[42]
7	c-lex.c:1213	last_insn→fld[1].rtx →fld[1].rtx→fld[3].rtx →fld[1].rtx.code
8	c-decl.c:1213	sequence_result[2] →fld[0].rtvec →elem[0].rtx→fld[1].rtx →fld[1].rtx→fld[1].rtx →fld[1].rtx→fld[1].rtx →fld[1].rtx→fld[1].rtx →fld[3].rtx→fld[1].rtx.code
9	combine.c:4271	x→fld[0].rtx→fld[0].rtx

combine.c:4279

```
if (GET_CODE (XEXP (x, 0)) == PLUS {  
  x = apply_distributive_law  
    (gen_binary (PLUS, mode,  
                gen_binary (MULT, mode,  
                            XEXP (XEXP (x, 0), 0),  
                            XEXP (x, 1)),  
                 gen_binary (MULT, mode,  
                            XEXP (XEXP (x, 0), 1),  
                            XEXP (x, 1))));  
  
  if (GET_CODE (x) != MULT)  
    return x;  
}
```

Should be copy_rtx()

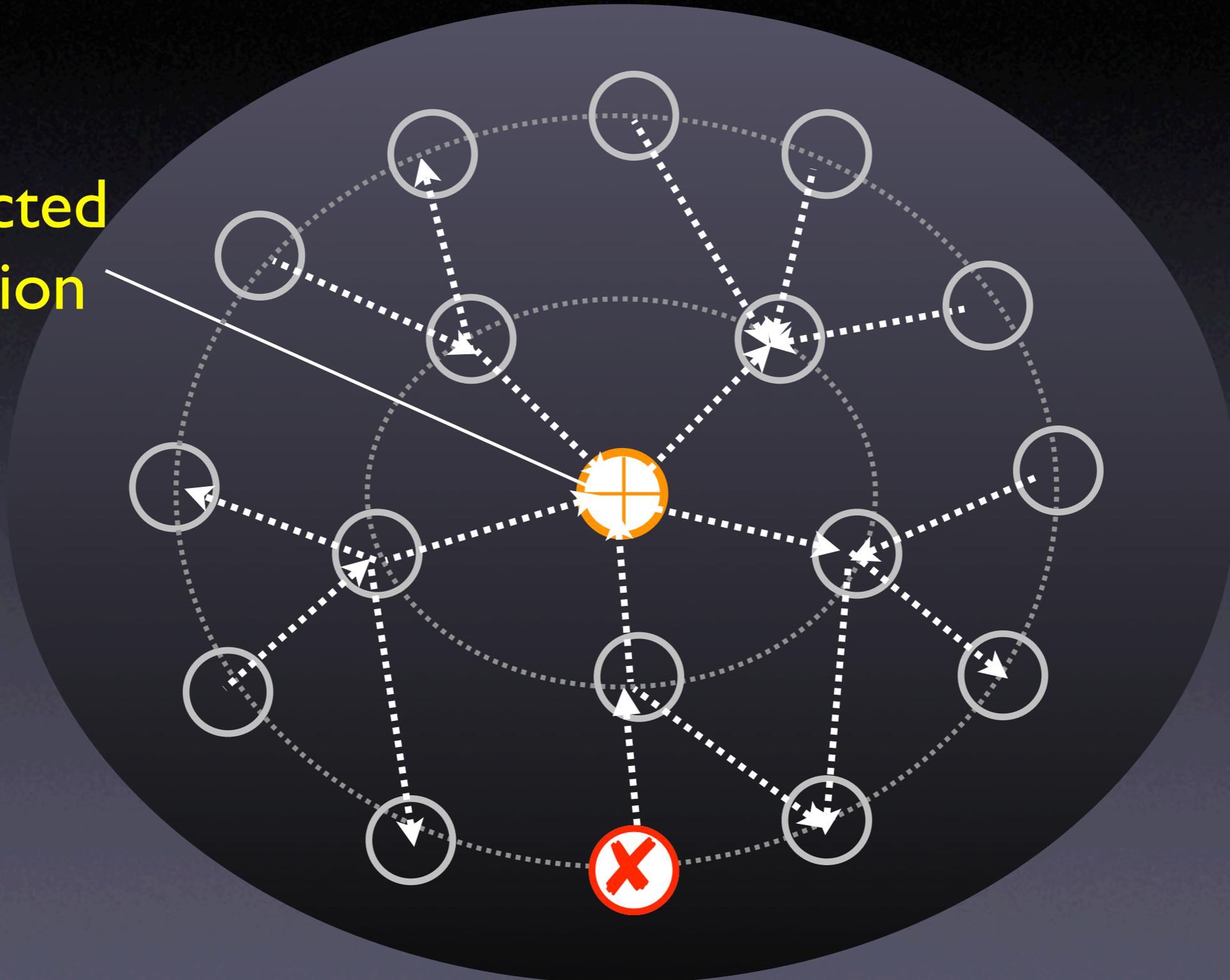
How good are we?

Evaluation using the **Siemens Testsuite**:

- 7 programs – most text processors
- 132 variations, each with 1 seeded defect
- Challenge: Using test runs, locate defect
- All proposed defect locators fail
(Comparing coverage, slicing, dynamic invariants)

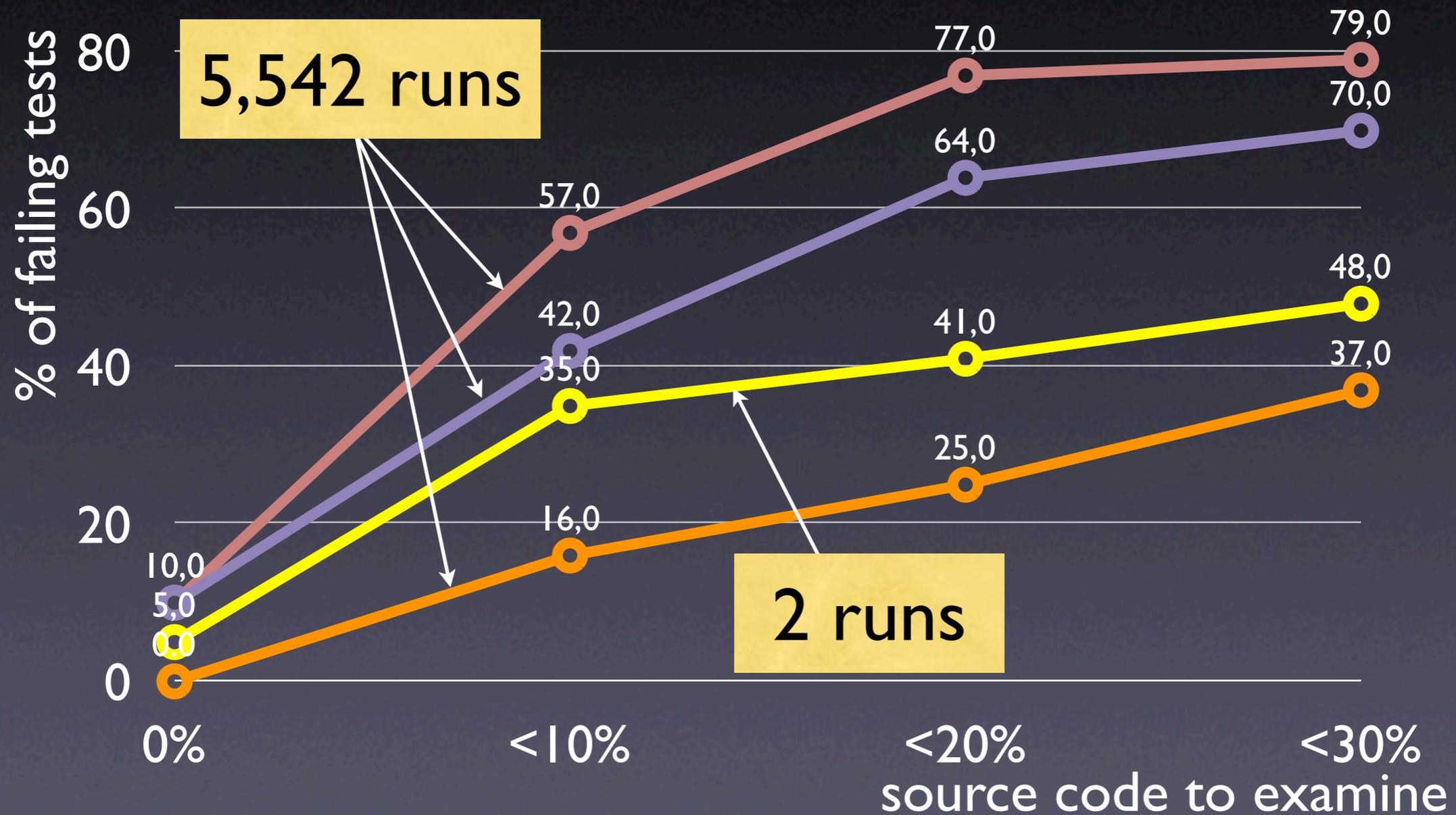
Close to the Defect

Predicted location



Locating Defects

- NN (Renieris + Reiss, ASE 2003)
- CT (Cleve + Zeller, ICSE 2005)
- SD (Liblit et al., PLDI 2005)
- SOBER (Liu et al, TR 2005)



Open Issues

- Hierarchical search
- Ranking transitions
- User-side diagnosis
- Combination with statistical causality

Concepts

- ★ Cause transitions pinpoint *failure causes in the program code*
- ★ Failure-causing statements are *potential fixes* (and frequently defects, too)
- ★ Even more demanding, yet effective technique

