

Mining Anomalies

Andrzej Wasylkowski

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Why Mine Anomalies?

- How can we make programs more reliable?
 - Testing, code inspection, etc.
 - Mining anomalies, etc.
 - In general: automatic defect detection

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Overview

Automatic Defect Detection

Rule-based Techniques

Specification-checking Techniques

Mining-based Techniques

Mining Repositories

Mining Traces

Mining Source Code

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Overview

Automatic Defect Detection

Rule-based Techniques

Specification-checking Techniques

Mining-based Techniques

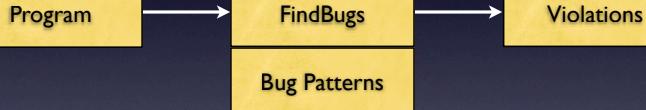
Mining Repositories

Mining Traces

Mining Source Code

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FindBugs



Hovemeyer, David, and William Pugh. 2004. Finding bugs is easy. *SIGPLAN Notices* 39, no. 12 (December): 92–106

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FindBugs's Bug Patterns

- Equal Objects Must Have Equal Hashcodes
- Static Field Modifiable By Untrusted Code
- Null Pointer Dereference
- Return Value Should Be Checked
- ...

Hovemeyer, David, and William Pugh. 2004. Finding bugs is easy. *SIGPLAN Notices* 39, no. 12 (December): 92–106

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Rule-based Techniques

- Fixed “bug patterns” to check against
- Pros: Fully automatic, scalable
- Cons: Limited to occurrences of “bug patterns”

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Rule-based Techniques

- Fixed “bug patterns” to check against
- **Can we add our own rules?**
- ~~Limited to occurrences of “bug patterns”~~

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Overview



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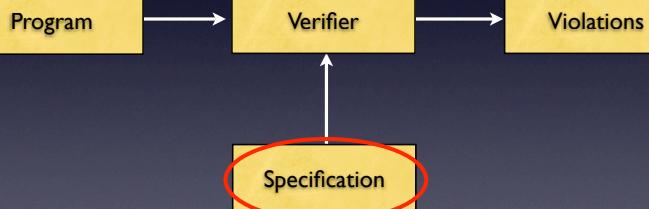
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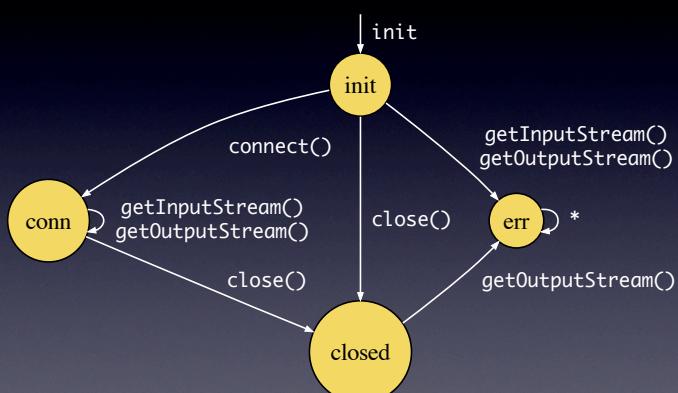
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Specification-checking



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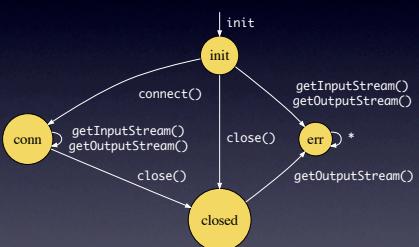
Typestate: java.net.Socket



Fink, Stephen J., Eran Yahav, Nurit Dor, G. Ramalingam, Emmanuel Geay.
2008. Effective typestate verification in the presence of aliasing. *ACM Transactions on Software Engineering and Methodology* 17, no. 2 (April): 1–34.

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Typestate Verification



```
...  
Socket s1 = new Socket();  
s1.connect(...);  
inp = s1.getInputStream();  
data = readData(inp);  
s1.close();  
...  
...  
Socket s1 = new Socket();  
inp = s1.getInputStream();  
data = readData(inp);  
s1.close();  
...
```

Fink, Stephen J., Eran Yahav, Nurit Dor, G. Ramalingam, Emmanuel Geay,
2008. Effective typestate verification in the presence of aliasing. ACM
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Specification-checking Techniques

- Use external specification to check against
- Pros: adaptable, very precise
- Cons: need specification, may have scalability problems

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Specification-checking Techniques

- Use external specification to check against
- Writing specifications is **very difficult!**
- Cons: need specification, may have scalability problems

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Mining-based Techniques

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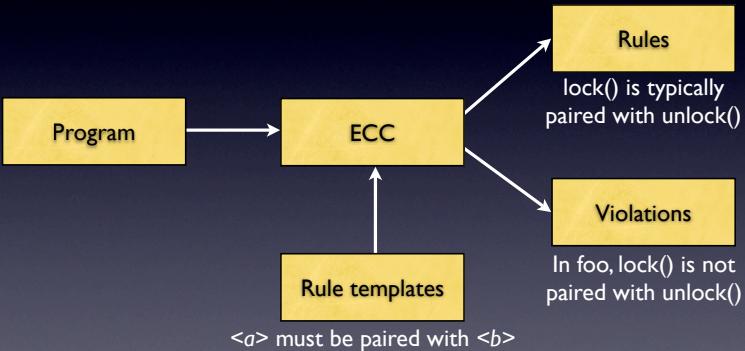
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Mining Source Code

- Code is **typically** correct
- **Deviant** behavior can point to a bug
- We can learn what is **common** behavior...
- ...and detect **uncommon** behavior

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ECC



Engler, Dawson, David Yu Chen, Seth Hallem, Andy Chou, and Benjamin Chelf. 2001. Bugs as deviant behavior: A general approach to inferring errors in systems code. In *SOSP 2001*, 57–72. New York, NY:ACM.

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ECC: Example

```

lock l;          // Lock
int a, b;        // Variables potentially
                 // protected by l
void foo () {
    lock (l);   // Enter critical section
    a = a + b;  // MAY: a,b protected by l
    unlock (l); // Exit critical section
    b = b + 1;  // MUST: b not protected by l
}
void bar () {
    lock (l);
    a = a + 1; // MAY: a protected by l
    unlock (l);
}
void baz () {
    a = a + 1; // MAY: a protected by l
    unlock (l);
    b = b - 1; // MUST: b not protected by l
    a = a / 5; // MUST: a not protected by l
}
  
```

Rule template:
lock </> protects variable <v>

Rule:
lock l protects variable a

Rule:
lock l protects variable b

Engler, Dawson, David Yu Chen, Seth Hallem, Andy Chou, and Benjamin Chelf. 2001. Bugs as deviant behavior: A general approach to inferring errors in systems code. In *SOSP 2001*, 57–72. New York, NY:ACM.

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Rule:
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Violation:
a is not protected by l in baz

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ECC: Summary

- Mines rules based on templates
- Pros: fully automatic, project-specific
- Cons: templates are simple and have fixed size

Engler, Dawson, David Yu Chen, Seth Hallem, Andy Chou, and Benjamin Cheff. 2001. Bugs as deviant behavior: A general approach to inferring errors in systems code. In SOSP 2001, 57–72. New York, NY:ACM.

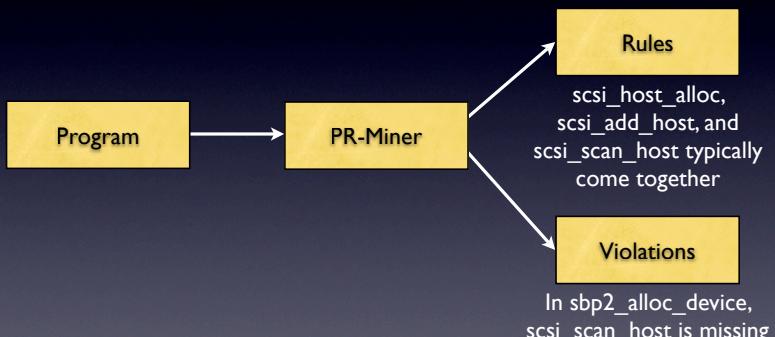
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ECC: Summary

- Mines rules based on templates
 - Templates have a **fixed** number of slots.
 - ~~Complex templates are simple and have fixed size~~

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PR-Miner



Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In *ESEC/FSE-13*, 306–315. New York, NY: ACM

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PR-Miner: Step I

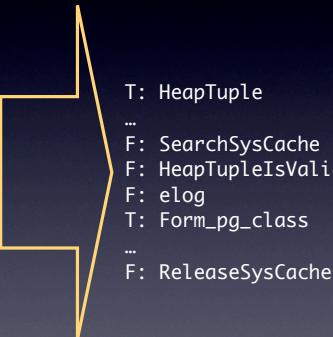
```
static void  
getRelationDescription (...)  
{  
    HeapTuple relTup;  
    ...  
    relTup = SearchSysCache (...);  
    if (!HeapTupleIsValid (relTup))  
        elog (...);  
    relForm = ...;  
    ...  
    ReleaseSysCache (relTup);  
}
```

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In *ESEC/FSE-13*, 306–315, New York, NY:ACM

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PR-Miner: Step 1

```
static void  
getRelationDescription (...)  
{  
    HeapTuple relTup;  
    ...  
    relTup = SearchSysCache (...);  
    if (!HeapTupleIsValid (relTup))  
        elog (...);  
    relForm = ...;  
    ...  
    ReleaseSysCache (relTup);  
}
```



T: HeapTuple
..
F: SearchSysCache
F: HeapTupleIsValid
F: elog
T: Form_pg_class
..
F: ReleaseSysCache

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In ESEC/FSE-13, 306–315, New York, NY:ACM

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PR-Miner: Step 2

T: HeapTuple	T: Form_pg_class
F: SearchSysCache	T: HeapTuple
F: HeapTupleIsValid	F: SearchSysCache
T: Form_pg_class	F: elog
F: ReleaseSysCache	F: ReleaseSysCache
...	...
T: StringInfoData	T: HeapTuple
T: HeapTuple	F: SearchSysCache
F: SearchSysCache	F: HeapTupleIsValid
F: NameStr	F: elog
F: ReleaseSysCache	T: Form_pg_class
...	...

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In ESEC/FSE-13, 306–315, New York, NY:ACM

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PR-Miner: Step 2

T: HeapTuple	T: Form_pg_class	Rule:
F: SearchSysCache	T: HeapTuple	T: HeapTuple,
F: HeapTupleIsValid	F: SearchSysCache	F: SearchSysCache, and
T: Form_pg_class	F: elog	F: ReleaseSysCache
F: ReleaseSysCache	F: ReleaseSysCache	typically come
...	...	together
T: StringInfoData	T: HeapTuple	
T: HeapTuple	F: SearchSysCache	
F: SearchSysCache	F: HeapTupleIsValid	
F: NameStr	F: elog	
F: ReleaseSysCache	T: Form_pg_class	
...		

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In ESEC/FSE-13, 306–315, New York, NY:ACM

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PR-Miner: Step 2

T: HeapTuple
F: SearchSysCache
F: HeapTupleIsValid
T: Form_pg_class
F: ReleaseSysCache
...

T: Form_pg_class
T: HeapTuple
F: SearchSysCache
F: elog
F: ReleaseSysCache
...

Rule:
T: HeapTuple,
F: SearchSysCache, and
F: ReleaseSysCache
typically come
together

T: StringInfoData
T: HeapTuple
F: SearchSysCache
F: NameStr
F: ReleaseSysCache
...

T: HeapTuple
F: SearchSysCache
F: HeapTupleIsValid
F: elog
T: Form_pg_class
...

Violation:
F: ReleaseSysCache is
missing

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In *ESEC/FSE-13*, 306–315, New York, NY:ACM

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PR-Miner: Summary

- Mines rules being sets of entities
- Pros: scalable, project-specific, flexible rule size
- Cons: no ordering of entities

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In *ESEC/FSE-13*, 306–315, New York, NY:ACM

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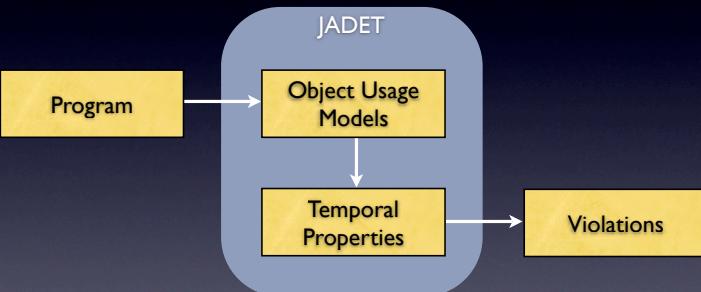
PR-Miner: Summary

- Mines rules being sets of entities
- Ordering is not taken into account.
- Cons: no ordering of entities

Li, Zhenmin, and Yuanyuan Zhou. 2005. PR-Miner: Automatically extracting implicit programming rules and detecting violations in large software code. In *ESEC/FSE-13*, 306–315, New York, NY:ACM

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JADET



Wasylkowski, Andrzej, Andreas Zeller, and Christian Lindig. 2007. Detecting object usage anomalies. In ESEC-FSE 2007, 35–44, New York, NY:ACM

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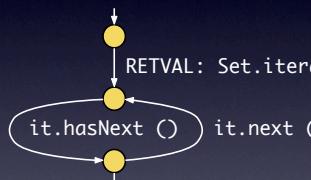
Creating an Object Usage Model: Example I

```
public List getList (Set ps) {  
    List l = new ArrayList ();  
    createList (this.cl, l);  
    Iterator it = ps.iterator ();  
    while (it.hasNext ()) {  
        Property p = it.next ();  
        addProperty (p, l);  
    }  
    reapList (l);  
    return l;  
}
```

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Creating an Object Usage Model: Example I

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public List getList (Set ps) {  
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```



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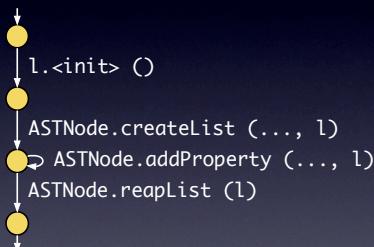
Creating an Object Usage Model: Example 2

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public List getList (Set ps) {  
    List l = new ArrayList ();  
    createList (this.cl, l);  
    Iterator it = ps.iterator ();  
    while (it.hasNext ()) {  
        Property p = it.next ();  
        addProperty (p, l);  
    }  
    reapList (l);  
    return l;  
}
```

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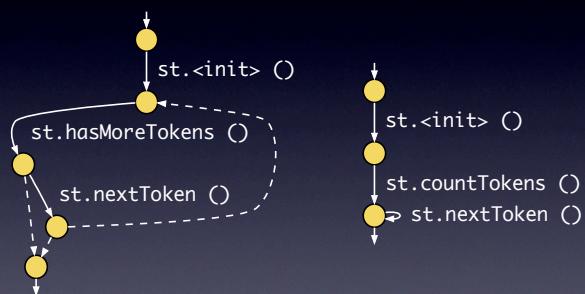
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```



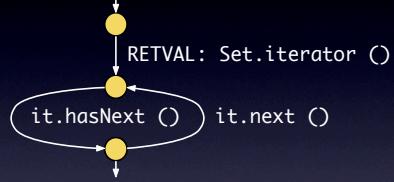
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Example OUMs: StringTokenizer



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Extracting Temporal Properties



Temporal properties

```
RETVAL: Set.iterator() < Iterator.hasNext() @ this
RETVAL: Set.iterator() < Iterator.next() @ this
Iterator.hasNext() @ this < Iterator.next() @ this
Iterator.hasNext() @ this < Iterator.hasNext() @ this
  Iterator.next() @ this < Iterator.hasNext() @ this
  Iterator.next() @ this < Iterator.next() @ this
```

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Extracting Temporal Properties

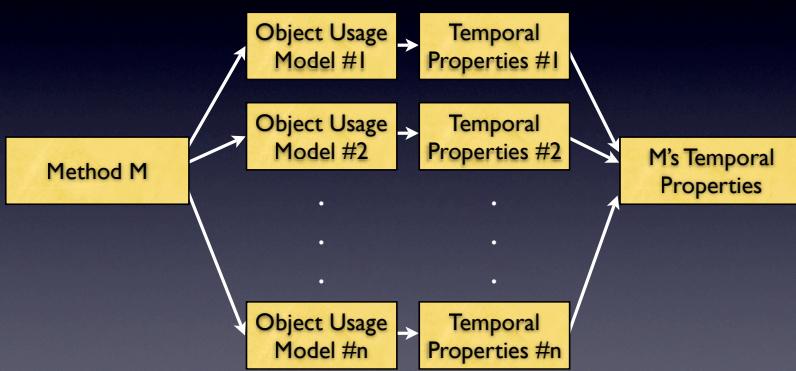


Temporal properties

```
StringTokenizer.<init>() @ this < StringTokenizer.countTokens() @ this
StringTokenizer.<init>() @ this < StringTokenizer.nextToken() @ this
StringTokenizer.countTokens() @ this < StringTokenizer.nextToken() @ this
StringTokenizer.nextToken() @ this < StringTokenizer.nextToken() @ this
```

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Extracting Temporal Properties: Summary



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Methods vs. Temporal Properties

	Temporal Properties				
	a<b	c<d	a<c	d<a	...
Methods	M1	■	■	■	
M2	■	■	■		
M3		■	■		⋮
M4		■	■	■	
...					

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Methods vs. Temporal Properties

	Temporal Properties				
	a<b	c<d	a<c	d<a	...
Methods	M1	■	■	■	
M2	■	■	■		
M3		■	■		⋮
M4		■	■	■	
...					

This forms
a pattern

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Methods vs. Temporal Properties

	Temporal Properties				
	a<b	c<d	a<c	d<a	...
Methods	M1	■	■	■	
M2	■	■	■		
M3		■	■		⋮
M4		■	■	■	
...					

Another
pattern

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Methods vs. Temporal Properties

	Temporal Properties				
	a<b	c<d	a<c	d<a	...
Methods	M1				
	M2				
	M3				
	M4				

Yet another pattern

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Detecting Violations

	Temporal Properties				
	a<b	c<d	a<c	d<a	...
Methods	M1				
	M2				
	M3				
	M4				

...

...

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Example Violation (I)

```
private boolean verifyNIAP (...) {  
    ...  
    Iterator iter = ...;  
    while (iter.hasNext()) {  
        ... = iter.next();  
        ...  
        return verifyNIAP (...);  
    }  
    return true;  
}
```

should be fixed

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Example Violation (2)

```
public String getRetentionPolicy () {  
    ...  
    for (Iterator it = ...; it.hasNext(); ) {  
        ... = it.next();  
        ...  
        return retentionPolicy;  
    }  
    ...  
}
```

should be fixed

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Example Violation (3)

```
public void visitCALOAD (CALOAD o) {  
    Type arrayref = stack().peek(1);  
    Type index = stack().peek(0);  
    indexOfInt(o, index);  
    arrayrefOfArrayType(o, arrayref);  
}  
  
should check the elements' type, too
```

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JADET: Summary

- Mines rules being sets of temporal properties
- Pros: fully automatic, scalable, project specific
- Cons: quite complicated, many false positives

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JADET: Summary

- Mines rules being sets of temporal properties
- All problems solved? Of course not!
- Cons: quite complicated,
many false positives

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Summary

- ★ Three main approaches:
 - ★ Rule-based Techniques
 - ★ Specification-checking techniques
 - ★ Mining-based techniques

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Code-mining Techniques

- ★ “Learn” rules from source code
- ★ Rule violation = potential defect
 - ★ Can find project-specific bugs
- ★ Many different rules types

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